

Pseudocodes der RB-Tree-Funktionen nach Cormen, Algorithms, MIT Press 2nd ed, 2001

RB-Insert(*T*, *z*) und *RB-Insert-Fixup*(*T*, *z*)

RB-Delete(*T*, *z*) und *RB-Delte-Fixup*(*T*, *x*)

RB-INSERT(*T*, *z*)

```
1  y ← nil[T]
2  x ← root[T]
3  while x ≠ nil[T]
4      do y ← x
5          if key[z] < key[x]
6              then x ← left[x]
7              else x ← right[x]
8  p[z] ← y
9  if y = nil[T]
10     then root[T] ← z
11     else if key[z] < key[y]
12         then left[y] ← z
13         else right[y] ← z
14 left[z] ← nil[T]
15 right[z] ← nil[T]
16 color[z] ← RED
17 RB-INSERT-FIXUP(T, z)
```

RB-INSERT-FIXUP(*T*, *z*)

```
1  while color[p[z]] = RED
2      do if p[z] = left[p[p[z]]]
3          then y ← right[p[p[z]]]
4              if color[y] = RED
5                  then color[p[z]] ← BLACK           ▷ Case 1
6                  color[y] ← BLACK                 ▷ Case 1
7                  color[p[p[z]]] ← RED             ▷ Case 1
8                  z ← p[p[z]]                       ▷ Case 1
9              else if z = right[p[z]]
10                 then z ← p[z]                     ▷ Case 2
11                 LEFT-ROTATE(T, z)                 ▷ Case 2
12                 color[p[z]] ← BLACK               ▷ Case 3
13                 color[p[p[z]]] ← RED             ▷ Case 3
14                 RIGHT-ROTATE(T, p[p[z]])        ▷ Case 3
15             else (same as then clause
16                 with “right” and “left” exchanged)
16  color[root[T]] ← BLACK
```

RB-DELETE(T, z)

```

1  if left[z] = nil[T] or right[z] = nil[T]
2      then y ← z
3      else y ← TREE-SUCCESSOR(z)
4  if left[y] ≠ nil[T]
5      then x ← left[y]
6      else x ← right[y]
7  p[x] ← p[y]
8  if p[y] = nil[T]
9      then root[T] ← x
10     else if y = left[p[y]]
11         then left[p[y]] ← x
12         else right[p[y]] ← x
13 if y ≠ z
14     then key[z] ← key[y]
15     copy y's satellite data into z
16 if color[y] = BLACK
17     then RB-DELETE-FIXUP(T, x)
18 return y

```

RB-DELETE-FIXUP(T, x)

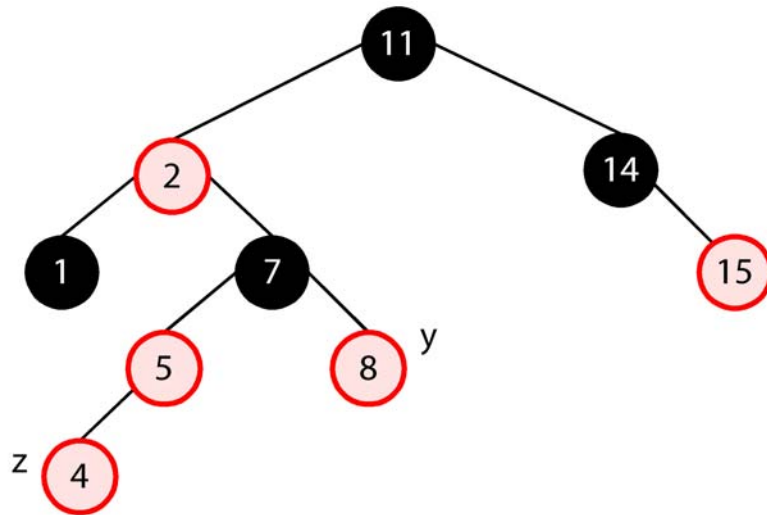
```

1  while x ≠ root[T] and color[x] = BLACK
2      do if x = left[p[x]]
3          then w ← right[p[x]]
4              if color[w] = RED
5                  then color[w] ← BLACK           ▷ Case 1
6                  color[p[x]] ← RED                 ▷ Case 1
7                  LEFT-ROTATE(T, p[x])              ▷ Case 1
8                  w ← right[p[x]]                   ▷ Case 1
9              if color[left[w]] = BLACK and color[right[w]] = BLACK
10                 then color[w] ← RED                ▷ Case 2
11                 x ← p[x]                            ▷ Case 2
12                 else if color[right[w]] = BLACK
13                     then color[left[w]] ← BLACK    ▷ Case 3
14                     color[w] ← RED                 ▷ Case 3
15                     RIGHT-ROTATE(T, w)             ▷ Case 3
16                     w ← right[p[x]]                ▷ Case 3
17                     color[w] ← color[p[x]]         ▷ Case 4
18                     color[p[x]] ← BLACK            ▷ Case 4
19                     color[right[w]] ← BLACK        ▷ Case 4
20                     LEFT-ROTATE(T, p[x])           ▷ Case 4
21                     x ← root[T]                   ▷ Case 4
22                 else (same as then clause with "right" and "left" exchanged)
23 color[x] ← BLACK

```

Die drei Fälle von RB-Insert-Fixup(T, z):

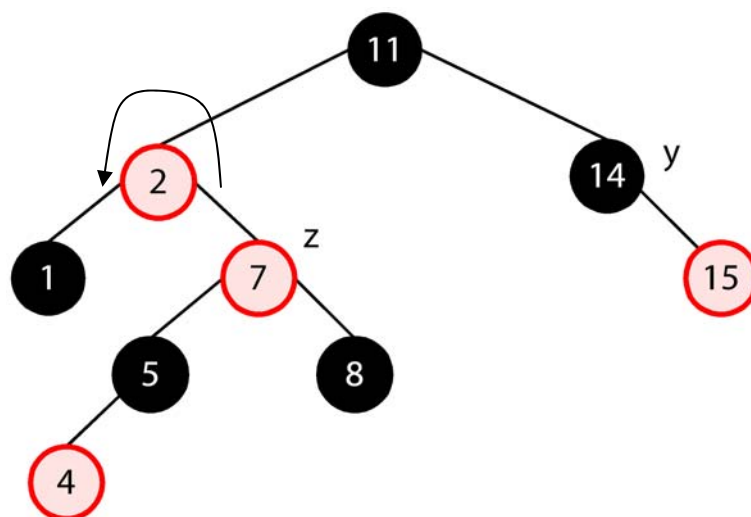
Fall 1: z 's Onkel y ist rot



→ Umfärben und z 's Großvater wird neues z .

Fall 2: z 's Onkel y ist schwarz und z ist rechtes Kind von $p[z]$

→ LINKS-ROTATION bei $p[z]$

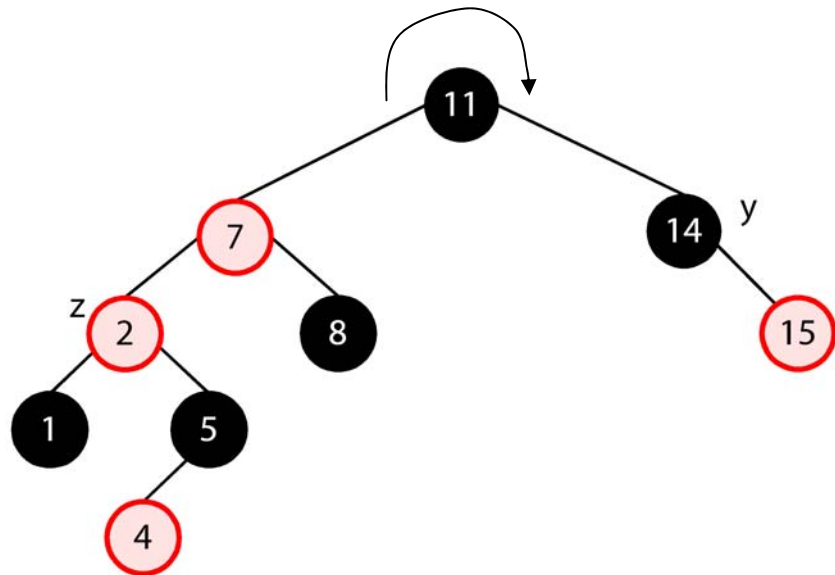


→ anschließende RECHTS-ROTATION

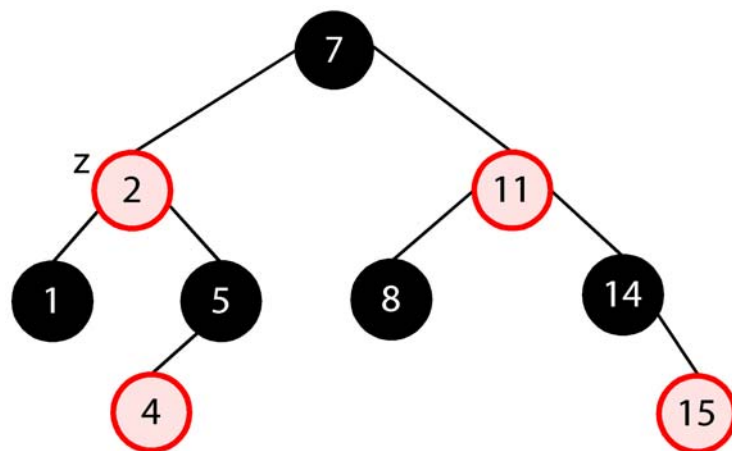
Beachte: Vor der LINKS-ROTATION wird $z \leftarrow p[z]$. Darum wird nach der LINKS-ROTATION z linkes Kind von $p[z]$. Automatischer Übergang zu Fall 3.

Fall 3: z's Onkel y ist schwarz und z ist linkes Kind von p[z]

→ RECHTS-ROTATION



→



Fertig!